

Weight enumerators of binary singly even self-dual codes

Akihiro Munemasa

Tohoku University

(joint work with Stefka Bouyuklieva and Masaaki Harada)

JCCA, August 18, 2017

Kumamoto University

Binary codes and their dual

A code C of length n is a vector subspace of \mathbb{F}_2^n .

The **dual** code C^\perp of C is defined as

$$C^\perp = \{x \in \mathbb{F}_2^n \mid x \cdot y = 0 \text{ for all } y \in C\},$$

and C is **self-dual** if $C = C^\perp$.

A self-dual code C is

doubly even $\iff \text{wt}(x) \equiv 0 \pmod{4} \quad (\forall x \in C),$

singly even \iff otherwise

$$\iff C_0 = \{x \in C \mid \text{wt}(x) \equiv 0 \pmod{4}\} \\ \subsetneq C \quad (\text{codimension } 1)$$

The **minimum weight** of C is

$$d(C) = \min\{\text{wt}(x) \mid 0 \neq x \in C\}.$$

Theorem (Mallows–Sloane (1973))

If C is a *doubly even* self-dual code of length n , then its minimum weight is at most

$$4 \left\lfloor \frac{n}{24} \right\rfloor + 4.$$

Upper bounds on minimum weight

Theorem (Mallows–Sloane (1973))

If C is a **doubly even** self-dual code of length n , then its minimum weight is at most

$$4 \left\lfloor \frac{n}{24} \right\rfloor + 4.$$

doubly even $\iff \text{wt}(x) \equiv 0 \pmod{4} \quad (\forall x \in C),$

singly even \iff otherwise

Upper bounds on minimum weight

Theorem (Mallows–Sloane (1973))

If C is a **doubly even** self-dual code of length n , then its minimum weight is at most

$$4 \left\lfloor \frac{n}{24} \right\rfloor + 4.$$

Theorem (Rains (1999))

If C is a self-dual code of length n , then its minimum weight is at most

$$\begin{cases} 4 \left\lfloor \frac{n}{24} \right\rfloor + 4 & \text{if } n \not\equiv 22 \pmod{24}, \\ 4 \left\lfloor \frac{n}{24} \right\rfloor + 6 & \text{if } n \equiv 22 \pmod{24}. \end{cases}$$

A self-dual code meeting this upper bound is called **extremal**.

Doubly even $\implies n \equiv 0 \pmod{8}$.

Shadow

Let C be a singly even self-dual code of length n and let

$$C_0 = \{x \in C \mid \text{wt}(x) \equiv 0 \pmod{4}\} \subsetneq C.$$

The **shadow** S is defined to be

$$S = C_0^\perp \setminus C.$$

Then

$$\text{wt}(x) \equiv n/2 \pmod{4} \quad (\forall x \in S),$$

so, letting $d(S) = \min\{\text{wt}(x) \mid x \in S\}$, we say that C is a code with **minimal shadow** if

$$d(S) = \begin{cases} 1 & \text{if } n \equiv 2 \pmod{8}, \\ 2 & \text{if } n \equiv 4 \pmod{8}, \\ 3 & \text{if } n \equiv 6 \pmod{8}, \\ 4 & \text{if } n \equiv 0 \pmod{8}. \end{cases}$$

Shadow

Let C be a singly even self-dual code of length n and let

$$C_0 = \{x \in C \mid \text{wt}(x) \equiv 0 \pmod{4}\} \subsetneq C.$$

The **shadow** S is defined to be

$$S = C_0^\perp \setminus C.$$

Then

$$\text{wt}(x) \equiv n/2 \pmod{4} \quad (\forall x \in S),$$

so, letting $d(S) = \min\{\text{wt}(x) \mid x \in S\}$, we say that C is a code with **minimal shadow** if

$$d(S) = \begin{cases} 1 & \text{if } n \equiv 2 \pmod{8}, \\ 2 & \text{if } n \equiv 4 \pmod{8}, \\ 3 & \text{if } n \equiv 6 \pmod{8}, \\ 4 & \text{if } n \equiv 0 \pmod{8}. \end{cases}$$

Shadow

Let C be a singly even self-dual code of length n and let

$$C_0 = \{x \in C \mid \text{wt}(x) \equiv 0 \pmod{4}\} \subsetneq C.$$

The **shadow** S is defined to be

$$S = C_0^\perp \setminus C.$$

Then

$$\text{wt}(x) \equiv n/2 \pmod{4} \quad (\forall x \in S),$$

so, letting $d(S) = \min\{\text{wt}(x) \mid x \in S\}$, we say that C is a code with **minimal shadow** if

$$d(S) = \begin{cases} 1 & \text{if } n \equiv 2 \pmod{8}, \\ 2 & \text{if } n \equiv 4 \pmod{8}, \\ 3 & \text{if } n \equiv 6 \pmod{8}, \\ 4 & \text{if } n \equiv 0 \pmod{8}. \end{cases}$$

Singly even self-dual codes with minimal shadow

length n	$d = 4m + 4$ (extremal)	$d = 4m + 2$
$24m + 2$	$d(C) = 4m + 4$, $\overline{\exists}$	$d(C) = 4m + 2$, !w.e. ($\overline{\exists}$)
$24m + 4$	$d(C) = 4m + 4$, $\overline{\exists}$	$d(C) = 4m + 2$, !w.e. ($\overline{\exists}$)
$24m + 6$	$d(C) = 4m + 4$, $\overline{\exists}$	$d(C) = 4m + 2$, !w.e.
$24m + 8$	$d(C) = 4m + 4$, !w.e. ($\overline{\exists}$)	
$24m + 10$	$d(C) = 4m + 4$, $\overline{\exists}$	$d(C) = 4m + 2$, !w.e. ($\overline{\exists}$)
$24m + 12$	$d(C) = 4m + 4$, !w.e. ($\overline{\exists}$)	
$24m + 14$	$d(C) = 4m + 4$, !w.e. ($\overline{\exists}$)	
$24m + 16$	$d(C) = 4m + 4$, ($\overline{\exists}$)	
$24m + 18$	$d(C) = 4m + 4$, !w.e. ($\overline{\exists}$)	
$24m + 22$	$d(C) = 4m + 6$, $\overline{\exists}$	$d(C) = 4m + 4$, !w.e.

($\overline{\exists}$) means that nonexistence is shown for sufficiently large m .

Bouyuklieva and Varbanov (2011), by Bouyuklieva and Willems (2012),

Bouyuklieva, Harada and Munemasa (2017+)

Singly even self-dual codes with minimal shadow

length n	$d = 4m + 4$ (extremal)	$d = 4m + 2$
$24m + 2$	$d(C) = 4m + 4, \bar{\exists}$	$d(C) = 4m + 2, \text{!w.e. } (\bar{\exists})$
$24m + 4$	$d(C) = 4m + 4, \bar{\exists}$	$d(C) = 4m + 2, \text{!w.e. } (\bar{\exists})$
$24m + 6$	$d(C) = 4m + 4, \bar{\exists}$	$d(C) = 4m + 2, \text{!w.e.}$
$24m + 8$	$d(C) = 4m + 4, \text{!w.e. } (\bar{\exists})$	
$24m + 10$	$d(C) = 4m + 4, \bar{\exists}$	$d(C) = 4m + 2, \text{!w.e. } (\bar{\exists})$
$24m + 12$	$d(C) = 4m + 4, \text{!w.e. } (\bar{\exists})$	
$24m + 14$	$d(C) = 4m + 4, \text{!w.e. } (\bar{\exists})$	
$24m + 16$	$d(C) = 4m + 4, (\bar{\exists})$	
$24m + 18$	$d(C) = 4m + 4, \text{!w.e. } (\bar{\exists})$	
$24m + 22$	$d(C) = 4m + 6, \bar{\exists}$	$d(C) = 4m + 4, \text{!w.e.}$

$(\bar{\exists})$ means that nonexistence is shown for sufficiently large m .

Bouyuklieva and Varbanov (2011), by Bouyuklieva and Willems (2012),

Bouyuklieva, Harada and Munemasa (2017+)

Singly even self-dual codes with minimal shadow

length n	$d = 4m + 4$ (extremal)	$d = 4m + 2$
$24m + 2$	$d(C) = 4m + 4, \bar{\exists}$	$d(C) = 4m + 2, !w.e. (\bar{\exists})$
$24m + 4$	$d(C) = 4m + 4, \bar{\exists}$	$d(C) = 4m + 2, !w.e. (\bar{\exists})$
$24m + 6$	$d(C) = 4m + 4, \bar{\exists}$	$d(C) = 4m + 2, !w.e.$
$24m + 8$	$d(C) = 4m + 4, !w.e. (\bar{\exists})$	
$24m + 10$	$d(C) = 4m + 4, \bar{\exists}$	$d(C) = 4m + 2, !w.e. (\bar{\exists})$
$24m + 12$	$d(C) = 4m + 4, !w.e. (\bar{\exists})$	
$24m + 14$	$d(C) = 4m + 4, !w.e. (\bar{\exists})$	
$24m + 16$	$d(C) = 4m + 4, (\bar{\exists})$	
$24m + 18$	$d(C) = 4m + 4, !w.e. (\bar{\exists})$	
$24m + 22$	$d(C) = 4m + 6, \bar{\exists}$	$d(C) = 4m + 4, !w.e.$

$(\bar{\exists})$ means that nonexistence is shown for sufficiently large m .

Bouyuklieva and Varbanov (2011), by Bouyuklieva and Willems (2012),

Bouyuklieva, Harada and Munemasa (2017+)

Singly even self-dual codes with minimal shadow

length n	$d = 4m + 4$ (extremal)	$d = 4m + 2$
$24m + 2$	$d(C) = 4m + 4, \bar{\exists}$	$d(C) = 4m + 2, \text{!w.e. } (\bar{\exists})$
$24m + 4$	$d(C) = 4m + 4, \bar{\exists}$	$d(C) = 4m + 2, \text{!w.e. } (\bar{\exists})$
$24m + 6$	$d(C) = 4m + 4, \bar{\exists}$	$d(C) = 4m + 2, \text{!w.e.}$
$24m + 8$	$d(C) = 4m + 4, \text{!w.e. } (\bar{\exists})$	
$24m + 10$	$d(C) = 4m + 4, \bar{\exists}$	$d(C) = 4m + 2, \text{!w.e. } (\bar{\exists})$
$24m + 12$	$d(C) = 4m + 4, \text{!w.e. } (\bar{\exists})$	
$24m + 14$	$d(C) = 4m + 4, \text{!w.e. } (\bar{\exists})$	
$24m + 16$	$d(C) = 4m + 4, (\bar{\exists})$	
$24m + 18$	$d(C) = 4m + 4, \text{!w.e. } (\bar{\exists})$	
$24m + 22$	$d(C) = 4m + 6, \bar{\exists}$	$d(C) = 4m + 4, \text{!w.e.}$

$(\bar{\exists})$ means that nonexistence is shown for sufficiently large m .

Bouyuklieva and Varbanov (2011), by Bouyuklieva and Willems (2012),

Bouyuklieva, Harada and Munemasa (2017+)

$$n = 24m + 4, d(C) = 4m + 2, d(S) = 2$$

The weight enumerators of C and S :

$$W_C(y) = \sum_{i=0}^{12m+2} a_i y^{2i} \equiv \mathbf{a} \mathbf{y}^\top \pmod{y^{6m+1}},$$

$$W_S(y) = \sum_{i=0}^{6m} b_i y^{4i+2} \equiv \mathbf{b} \mathbf{y}'^\top \pmod{y^{12m+3}},$$

where

$$\mathbf{y} = (1, y^2, y^4, \dots, y^{6m}) \in \mathbb{Q}[y]^{3m+1},$$

$$\mathbf{y}' = (y^2, y^6, y^{10}, \dots, y^{12m+2}) \in \mathbb{Q}[y]^{3m+1},$$

$$\mathbf{a} = (a_0, a_1, \dots, a_{3m}) \in \mathbb{Z}^{3m+1},$$

$$\mathbf{b} = (b_0, b_1, \dots, b_{3m}) \in \mathbb{Z}^{3m+1}.$$

$$n = 24m + 4, d(C) = 4m + 2, d(S) = 2$$

The weight enumerators of C and S :

$$W_C(y) = \sum_{i=0}^{12m+2} a_i y^{2i} \equiv \mathbf{a} \mathbf{y}^\top \pmod{y^{6m+1}},$$

$$W_S(y) = \sum_{i=0}^{6m} b_i y^{4i+2} \equiv \mathbf{b} \mathbf{y}'^\top \pmod{y^{12m+3}},$$

where

$$\mathbf{y} = (1, y^2, y^4, \dots, y^{6m}) \in \mathbb{Q}[y]^{3m+1},$$

$$\mathbf{y}' = (y^2, y^6, y^{10}, \dots, y^{12m+2}) \in \mathbb{Q}[y]^{3m+1},$$

$$\mathbf{a} = (a_0, a_1, \dots, a_{3m}) \in \mathbb{Z}^{3m+1},$$

$$\mathbf{b} = (b_0, b_1, \dots, b_{3m}) \in \mathbb{Z}^{3m+1}.$$

$$n = 24m + 4, d(C) = 4m + 2, d(S) = 2$$

The weight enumerators of C and S :

$$W_C(y) = \sum_{i=0}^{12m+2} a_i y^{2i} \equiv \mathbf{a} \mathbf{y}^\top \pmod{y^{6m+1}},$$

$$W_S(y) = \sum_{i=0}^{6m} b_i y^{4i+2} \equiv \mathbf{b} \mathbf{y}'^\top \pmod{y^{12m+3}},$$

where

$$\mathbf{y} = (1, y^2, y^4, \dots, y^{6m}) \in \mathbb{Q}[y]^{3m+1},$$

$$\mathbf{y}' = (y^2, y^6, y^{10}, \dots, y^{12m+2}) \in \mathbb{Q}[y]^{3m+1},$$

$$\mathbf{a} = (a_0, a_1, \dots, a_{3m}) \in \mathbb{Z}^{3m+1},$$

$$\mathbf{b} = (b_0, b_1, \dots, b_{3m}) \in \mathbb{Z}^{3m+1}.$$

$$\mathbf{y} = (1, y^2, y^4, \dots, y^{6m}),$$

$$\mathbf{y}' = (y^2, y^6, y^{10}, \dots, y^{12m+2})$$

Rains (1999): Let

$$f_j = (1 + y^2)^{12m+2-4j} (y^2(1 - y^2)^2)^j \in \mathbb{Q}[y^2],$$

$$g_j = (-1)^j 2^{12m+2-6j} y^{12m+2-4j} (1 - y^4)^{2j} \in y^2 \mathbb{Q}[y^4].$$

Then

$$\mathbb{Q}[y]^{3m+1} \ni \mathbf{f} = (f_0, f_1, \dots, f_{3m+1}), \quad \mathbf{y} \equiv \mathbf{f}A \pmod{y^{6m+1}},$$

$$\ni \mathbf{g} = (g_0, g_1, \dots, g_{3m+1}), \quad \mathbf{y}' \equiv \mathbf{g}B \pmod{y^{12m+3}}.$$

$\exists \mathbf{c} \in \mathbb{Q}^{3m+1}$ such that

$$a\mathbf{y}^\top \equiv W_C(\mathbf{y}) = \mathbf{c}\mathbf{f}^\top \implies a\mathbf{A}^\top = \mathbf{c}$$

$$b\mathbf{y}'^\top \equiv W_S(\mathbf{y}) = \mathbf{c}\mathbf{g}^\top \implies b\mathbf{B}^\top = \mathbf{c}$$

$$\mathbf{y} = (1, y^2, y^4, \dots, y^{6m}),$$

$$\mathbf{y}' = (y^2, y^6, y^{10}, \dots, y^{12m+2})$$

Rains (1999): Let

$$f_j = (1 + y^2)^{12m+2-4j} (y^2(1 - y^2)^2)^j \in \mathbb{Q}[y^2],$$

$$g_j = (-1)^j 2^{12m+2-6j} y^{12m+2-4j} (1 - y^4)^{2j} \in y^2 \mathbb{Q}[y^4].$$

Then

$$\mathbb{Q}[y]^{3m+1} \ni \mathbf{f} = (f_0, f_1, \dots, f_{3m+1}), \quad \mathbf{y} \equiv \mathbf{f}A \pmod{y^{6m+1}},$$

$$\ni \mathbf{g} = (g_0, g_1, \dots, g_{3m+1}), \quad \mathbf{y}' \equiv \mathbf{g}B \pmod{y^{12m+3}}.$$

$\exists \mathbf{c} \in \mathbb{Q}^{3m+1}$ such that

$$a\mathbf{y}^\top \equiv W_C(\mathbf{y}) = \mathbf{c}\mathbf{f}^\top \implies a\mathbf{A}^\top = \mathbf{c}$$

$$b\mathbf{y}'^\top \equiv W_S(\mathbf{y}) = \mathbf{c}\mathbf{g}^\top \implies b\mathbf{B}^\top = \mathbf{c}$$

$$\mathbf{y} = (1, y^2, y^4, \dots, y^{6m}),$$

$$\mathbf{y}' = (y^2, y^6, y^{10}, \dots, y^{12m+2})$$

Rains (1999): Let

$$f_j = (1 + y^2)^{12m+2-4j} (y^2(1 - y^2)^2)^j \in \mathbb{Q}[y^2],$$

$$g_j = (-1)^j 2^{12m+2-6j} y^{12m+2-4j} (1 - y^4)^{2j} \in y^2 \mathbb{Q}[y^4].$$

Then

$$\mathbb{Q}[y]^{3m+1} \ni \mathbf{f} = (f_0, f_1, \dots, f_{3m+1}), \quad \mathbf{y} \equiv \mathbf{f}A \pmod{y^{6m+1}},$$

$$\ni \mathbf{g} = (g_0, g_1, \dots, g_{3m+1}), \quad \mathbf{y}' \equiv \mathbf{g}B \pmod{y^{12m+3}}.$$

$\exists \mathbf{c} \in \mathbb{Q}^{3m+1}$ such that

$$\mathbf{a}\mathbf{y}^\top \equiv W_C(\mathbf{y}) = \mathbf{c}\mathbf{f}^\top \implies \mathbf{a}A^\top = \mathbf{c}$$

$$\mathbf{b}\mathbf{y}'^\top \equiv W_S(\mathbf{y}) = \mathbf{c}\mathbf{g}^\top \implies \mathbf{b}B^\top = \mathbf{c}$$

$$\mathbf{y} = (1, y^2, y^4, \dots, y^{6m}),$$

$$\mathbf{y}' = (y^2, y^6, y^{10}, \dots, y^{12m+2})$$

Rains (1999): Let

$$f_j = (1 + y^2)^{12m+2-4j} (y^2(1 - y^2)^2)^j \in \mathbb{Q}[y^2],$$

$$g_j = (-1)^j 2^{12m+2-6j} y^{12m+2-4j} (1 - y^4)^{2j} \in y^2 \mathbb{Q}[y^4].$$

Then

$$\mathbb{Q}[y]^{3m+1} \ni \mathbf{f} = (f_0, f_1, \dots, f_{3m+1}), \quad \mathbf{y} \equiv \mathbf{f}A \pmod{y^{6m+1}},$$

$$\ni \mathbf{g} = (g_0, g_1, \dots, g_{3m+1}), \quad \mathbf{y}' \equiv \mathbf{g}B \pmod{y^{12m+3}}.$$

$\exists \mathbf{c} \in \mathbb{Q}^{3m+1}$ such that

$$\mathbf{a}\mathbf{y}^\top \equiv W_C(\mathbf{y}) = \mathbf{c}\mathbf{f}^\top \implies \mathbf{a}A^\top = \mathbf{c}$$

$$\mathbf{b}\mathbf{y}'^\top \equiv W_S(\mathbf{y}) = \mathbf{c}\mathbf{g}^\top \implies \mathbf{b}B^\top = \mathbf{c}$$

$$\mathbf{y} = (1, y^2, y^4, \dots, y^{6m}),$$

$$\mathbf{y}' = (y^2, y^6, y^{10}, \dots, y^{12m+2})$$

Rains (1999): Let

$$f_j = (1 + y^2)^{12m+2-4j} (y^2(1 - y^2)^2)^j \in \mathbb{Q}[y^2],$$

$$g_j = (-1)^j 2^{12m+2-6j} y^{12m+2-4j} (1 - y^4)^{2j} \in y^2 \mathbb{Q}[y^4].$$

Then

$$\mathbb{Q}[y]^{3m+1} \ni \mathbf{f} = (f_0, f_1, \dots, f_{3m+1}), \quad \mathbf{y} \equiv \mathbf{f}A \pmod{y^{6m+1}},$$

$$\ni \mathbf{g} = (g_0, g_1, \dots, g_{3m+1}), \quad \mathbf{y}' \equiv \mathbf{g}B \pmod{y^{12m+3}}.$$

$\exists \mathbf{c} \in \mathbb{Q}^{3m+1}$ such that

$$\mathbf{a}\mathbf{y}^\top \equiv W_C(\mathbf{y}) = \mathbf{c}\mathbf{f}^\top \implies \mathbf{a}A^\top = \mathbf{c}$$

$$\mathbf{b}\mathbf{y}'^\top \equiv W_S(\mathbf{y}) = \mathbf{c}\mathbf{g}^\top \implies \mathbf{b}B^\top = \mathbf{c}$$

$$\mathbf{a}A^\top = \mathbf{b}B^\top, \quad d(C) = 4m + 2, \quad d(S) = 2$$

$$W_C(y) = \sum_{i=0}^{12m+2} a_i y^{2i} \equiv \mathbf{a}\mathbf{y}^\top \pmod{y^{6m+1}}$$

$$d(C) = 4m + 2 \implies a_0 = 1, \quad a_1 = \cdots = a_{2m} = 0,$$

$$W_S(y) = \sum_{i=0}^{6m} b_i y^{4i+2} \equiv \mathbf{b}\mathbf{y}'^\top \pmod{y^{12m+3}}$$

$$d(S) = 2 \implies b_0 = 1, \quad b_1 = \cdots = b_{m-1} = 0.$$

$\mathbf{a}A^\top = \mathbf{b}B^\top$ implies

$$(1, \underbrace{0, \dots, 0}_{2m}, \underbrace{\mathbf{a}'}_m) \begin{pmatrix} * & * \\ 0 & A' \end{pmatrix} = (1, \underbrace{0, \dots, 0}_{m-1}, *) \begin{pmatrix} * & B' \\ * & 0 \end{pmatrix}$$

$$\implies \mathbf{a}'A' \rightarrow \mathbf{a}' \rightarrow \mathbf{a} \rightarrow W_C(y).$$

$$\mathbf{a}A^\top = \mathbf{b}B^\top, \quad d(C) = 4m + 2, \quad d(S) = 2$$

$$W_C(y) = \sum_{i=0}^{12m+2} a_i y^{2i} \equiv \mathbf{a}\mathbf{y}^\top \pmod{y^{6m+1}}$$

$$d(C) = 4m + 2 \implies a_0 = 1, \quad a_1 = \dots = a_{2m} = 0,$$

$$W_S(y) = \sum_{i=0}^{6m} b_i y^{4i+2} \equiv \mathbf{b}\mathbf{y}'^\top \pmod{y^{12m+3}}$$

$$d(S) = 2 \implies b_0 = 1, \quad b_1 = \dots = b_{m-1} = 0.$$

$\mathbf{a}A^\top = \mathbf{b}B^\top$ implies

$$(1, \underbrace{0, \dots, 0}_{2m}, \underbrace{\mathbf{a}'}_m) \begin{pmatrix} * & * \\ 0 & A' \end{pmatrix} = (1, \underbrace{0, \dots, 0}_{m-1}, *) \begin{pmatrix} * & B' \\ * & 0 \end{pmatrix}$$

$$\implies \mathbf{a}'A' \rightarrow \mathbf{a}' \rightarrow \mathbf{a} \rightarrow W_C(y).$$

$$\mathbf{a}A^\top = \mathbf{b}B^\top, \quad d(C) = 4m + 2, \quad d(S) = 2$$

$$W_C(y) = \sum_{i=0}^{12m+2} a_i y^{2i} \equiv \mathbf{a}\mathbf{y}^\top \pmod{y^{6m+1}}$$

$$d(C) = 4m + 2 \implies a_0 = 1, \quad a_1 = \cdots = a_{2m} = 0,$$

$$W_S(y) = \sum_{i=0}^{6m} b_i y^{4i+2} \equiv \mathbf{b}\mathbf{y}'^\top \pmod{y^{12m+3}}$$

$$d(S) = 2 \implies b_0 = 1, \quad b_1 = \cdots = b_{m-1} = 0.$$

$\mathbf{a}A^\top = \mathbf{b}B^\top$ implies

$$(1, \underbrace{0, \dots, 0}_{2m}, \underbrace{\mathbf{a}'}_m) \begin{pmatrix} * & * \\ 0 & A' \end{pmatrix} = (1, \underbrace{0, \dots, 0}_{m-1}, *) \begin{pmatrix} * & B' \\ * & 0 \end{pmatrix}$$

$$\implies \mathbf{a}'A' \rightarrow \mathbf{a}' \rightarrow \mathbf{a} \rightarrow W_C(y).$$

Formula for b_m

$$W_S(y) = \sum_{i=0}^{6m} b_i y^{4i+2} \quad (b_0 = 1, b_1 = \dots = b_{m-1} = 0)$$

is also uniquely determined. In fact,

$$b_m = \frac{2(12m+1)(38m+7)}{5m(2m+1)} \binom{5m}{m-1},$$

but incorrectly reported in Zhang–Michel–Feng–Ge (2015).

Moreover,

$$b_{m+1} = - \frac{\text{polynomial in } m \text{ of positive leading coeff.}}{(5m-1) \prod_{i=2}^6 (4m+i)} \binom{5m}{m-1} < 0 \quad (\text{for } m \text{ sufficiently large}).$$

Formula for b_m

$$W_S(y) = \sum_{i=0}^{6m} b_i y^{4i+2} \quad (b_0 = 1, b_1 = \dots = b_{m-1} = 0)$$

is also uniquely determined. In fact,

$$b_m = \frac{2(12m+1)(38m+7)}{5m(2m+1)} \binom{5m}{m-1},$$

but incorrectly reported in Zhang–Michel–Feng–Ge (2015).

Moreover,

$$b_{m+1} = - \frac{\text{polynomial in } m \text{ of positive leading coeff.}}{(5m-1) \prod_{i=2}^6 (4m+i)} \binom{5m}{m-1} < 0 \quad (\text{for } m \text{ sufficiently large}).$$

Formula for b_m

$$W_S(y) = \sum_{i=0}^{6m} b_i y^{4i+2} \quad (b_0 = 1, b_1 = \dots = b_{m-1} = 0)$$

is also uniquely determined. In fact,

$$b_m = \frac{2(12m+1)(38m+7)}{5m(2m+1)} \binom{5m}{m-1},$$

but incorrectly reported in Zhang–Michel–Feng–Ge (2015).

Moreover,

$$b_{m+1} = - \frac{\text{polynomial in } m \text{ of positive leading coeff.}}{(5m-1) \prod_{i=2}^6 (4m+i)} \binom{5m}{m-1} < 0 \quad (\text{for } m \text{ sufficiently large}).$$

Theorem (Bouyuklieva–Harada–M., arXiv:1707.04059)

The weight enumerators $W_C(y)$ and $W_S(y)$ of a singly even self-dual code C of length $24m + 4$, minimum weight $4m + 2$ and its shadow are uniquely determined by m . These uniquely determined polynomials have all coefficients nonnegative if and only if $0 \leq m \leq 155$.

In particular, for $m \geq 156$, there is no singly even self-dual code of length $24m + 4$, minimum weight $4m + 2$ with minimal shadow.

We have similar theorems for the lengths $24m + 2$ and $24m + 10$.

Theorem (Bouyuklieva–Harada–M., arXiv:1707.04059)

The weight enumerators $W_C(y)$ and $W_S(y)$ of a singly even self-dual code C of length $24m + 4$, minimum weight $4m + 2$ and its shadow are uniquely determined by m . These uniquely determined polynomials have all coefficients nonnegative if and only if $0 \leq m \leq 155$.

In particular, for $m \geq 156$, there is no singly even self-dual code of length $24m + 4$, minimum weight $4m + 2$ with minimal shadow.

We have similar theorems for the lengths $24m + 2$ and $24m + 10$.